DS 1300 - Introduction to SQL
Part 3 - Aggregation & other Topics

by Michael Hahsler
Based on slides for CS145 Introduction to Databases (Stanford)

DILBERT

I ACCOMPLISHED NOTHING THIS WEEK BECAUSE I WAS IN A TRAINING CLASS.

I DIDN'T APPROVE ANY TRAINING EXPENSES.

A VENDOR PAID FOR IT.

YOU DIDN'T ASK FOR PERMISSION.

I'M PRO-ACTIVE AND EMPOWERED.

AND WHAT WAS THE NAME OF THIS ALLEGED CLASS?

ADVANCED SCRIPTING STRUCTURE FOR INTERNETWORK OPTIMIZATION OF SQL DATABASES.

THAT DOESN'T SOUND REAL.

I CAN'T DO MY JOB IF YOU DON'T TRUST ME!

DO YOU LIKE HOW I COMBINED AGGRESSIVENESS WITH MY BASELINE LEVEL OF USELESSNESS?

I HAVE A GOOD FEELING ABOUT THIS.

YOU MIGHT NEED MORE AGGRESSIVENESS.
Lecture Overview

1. Aggregation & GROUP BY

2. Advanced SQL-izing (set operations, NULL, Outer Joins, etc.)
AGGREGATION, GROUP BY AND HAVING CLAUSE
Aggregation

SELECT COUNT(*)
FROM Product
WHERE year > 1995

SELECT AVG(price)
FROM Product
WHERE maker = 'Toyota'

• SQL supports several aggregation operations:
  • SUM, COUNT, MIN, MAX, AVG

Except for COUNT, all aggregations apply to a single attribute!
Aggregation: COUNT

COUNT counts the number of tuples including duplicates.

SELECT COUNT(category) 
FROM Product 
WHERE year > 1995

Note: Same as COUNT(*)!

We probably want count the number of “different” categories:

SELECT COUNT(DISTINCT category) 
FROM Product 
WHERE year > 1995
More Examples

Purchase(product, date, price, quantity)

SELECT SUM(price * quantity)
FROM Purchase

SELECT SUM(price * quantity)
FROM Purchase
WHERE product = 'bagel'

What do these mean?
Simple Aggregations

Purchase

<table>
<thead>
<tr>
<th>Product</th>
<th>Date</th>
<th>Price</th>
<th>Quantity</th>
</tr>
</thead>
<tbody>
<tr>
<td>bagel</td>
<td>10/21</td>
<td>1</td>
<td>20</td>
</tr>
<tr>
<td>banana</td>
<td>10/3</td>
<td>0.5</td>
<td>10</td>
</tr>
<tr>
<td>banana</td>
<td>10/10</td>
<td>1</td>
<td>10</td>
</tr>
<tr>
<td>bagel</td>
<td>10/25</td>
<td>1.50</td>
<td>20</td>
</tr>
</tbody>
</table>

SELECT SUM(price * quantity) FROM Purchase WHERE product = 'bagel'

50 (= 1*20 + 1.50*20)
Grouping and Aggregation

Find total sales after Oct 1, 2010, per product.

Let’s see what this means...

Note: Be very careful with dates! Use date/time related functions!
Grouping and Aggregation

Semantics of the query:

1. Compute the **FROM** and **WHERE** clauses

2. Group by the attributes in the **GROUP BY**

3. Compute the **SELECT** clause: grouped attributes and aggregates
1. Compute the **FROM** and **WHERE** clauses

```sql
SELECT product, SUM(price*quantity) AS TotalSales
FROM Purchase
WHERE date > '2000-10-01'
GROUP BY product
```

<table>
<thead>
<tr>
<th>Product</th>
<th>Date</th>
<th>Price</th>
<th>Quantity</th>
</tr>
</thead>
<tbody>
<tr>
<td>Bagel</td>
<td>2000-10-21</td>
<td>1</td>
<td>20</td>
</tr>
<tr>
<td>Bagel</td>
<td>2000-10-25</td>
<td>1.50</td>
<td>20</td>
</tr>
<tr>
<td>Banana</td>
<td>2000-10-03</td>
<td>0.5</td>
<td>10</td>
</tr>
<tr>
<td>Banana</td>
<td>2000-10-10</td>
<td>1</td>
<td>10</td>
</tr>
</tbody>
</table>
2. Group by the attributes in the \textbf{GROUP BY}

\begin{verbatim}
SELECT   product, SUM(price*quantity) AS TotalSales
FROM     Purchase
WHERE    date > '2000-10-01'
\end{verbatim}

\textbf{GROUP BY} \textit{product}
3. Compute the **SELECT** clause: grouped attributes and aggregates

```
SELECT product, SUM(price*quantity) AS TotalSales
FROM Purchase
WHERE date > '2000-10-01'
GROUP BY product
```

<table>
<thead>
<tr>
<th>Product</th>
<th>Date</th>
<th>Price</th>
<th>Quantity</th>
</tr>
</thead>
<tbody>
<tr>
<td>Bagel</td>
<td>2000-10-21</td>
<td>1</td>
<td>20</td>
</tr>
<tr>
<td></td>
<td>2000-10-25</td>
<td>1.50</td>
<td>20</td>
</tr>
<tr>
<td>Banana</td>
<td>2000-10-03</td>
<td>0.5</td>
<td>10</td>
</tr>
<tr>
<td></td>
<td>2000-10-10</td>
<td>1</td>
<td>10</td>
</tr>
</tbody>
</table>

**SELECT**

<table>
<thead>
<tr>
<th>Product</th>
<th>TotalSales</th>
</tr>
</thead>
<tbody>
<tr>
<td>Bagel</td>
<td>50</td>
</tr>
<tr>
<td>Banana</td>
<td>15</td>
</tr>
</tbody>
</table>
Activity

1) What do the next two queries calculate?

```sql
SELECT SUM(price) AS total, SUM(price) * 1.08 AS totalPlusTax
FROM Product pr
JOIN Purchase p ON pr.PName = p.product
WHERE p.buyer = 'Joe Blow'
```

```sql
SELECT p.buyer, SUM(price) AS total, SUM(price) * 1.08 AS totalPlusTax
FROM Product pr
JOIN Purchase p ON pr.PName = p.product
GROUP BY p.buyer
ORDER BY 1
```

2) Write a query to find the price of the most expensive product in each category.
Having Clause

Same query as before, except that we consider only products that have more than 100 buyers

```
SELECT product, SUM(price*quantity)
FROM Purchase
WHERE date > '2005-10-01'
GROUP BY product
HAVING SUM(quantity) > 100
```

Having clauses contain conditions on aggregates

Whereas where clauses condition on individual tuples...
General form of Grouping and Aggregation

```
SELECT S
FROM R_1, ..., R_n
WHERE C_1
GROUP BY a_1, ..., a_k
HAVING C_2
```

- \( S = \) Can ONLY contain attributes \( a_1, ..., a_k \) and/or aggregates over other attributes
- \( C_1 = \) is any condition on the attributes in \( R_1, ..., R_n \)
- \( C_2 = \) is any condition on the aggregate expressions
### General form of Grouping and Aggregation

```
SELECT S
FROM R_1, ..., R_n
WHERE C_1
GROUP BY a_1, ..., a_k
HAVING C_2
```

#### Evaluation steps:

1. **Evaluate FROM-WHERE**: apply condition $C_1$ on the attributes in $R_1, ..., R_n$
2. **GROUP BY** the attributes $a_1, ..., a_k$
3. Compute aggregates in $S$ and do projection (SELECT)
4. Apply condition $C_2$ to each group (may have aggregates)
Activity

1) What does this query do?

```
SELECT p.buyer, SUM(price) AS total, COUNT(*) AS purchases
FROM Product pr
JOIN Purchase p ON pr.PName = p.product
GROUP BY p.buyer
HAVING purchases >2
ORDER BY 1
```

2) What products in the DB have a revenue of more then $10,000?
OTHER SQL TOPICS: SUBQUERIES, NULLS, CASTING, OUTER JOINS AND ADDING DATA
Subqueries

```
SELECT *
FROM (SELECT product, COUNT(product) AS count
     FROM Purchase GROUP BY product)
WHERE count > 2
```

```
SELECT *, (SELECT count(*) FROM Product p1
           WHERE p1.category = p2.category) AS '# Prod. in Cat.'
FROM Product p2
```

Subqueries can appear wherever a table or a value is needed.
NULL VALUES & OTHER DETAILS
NULL Values

• Whenever we do not have a value, we can use NULL

• Can mean many things:
  – Value does not exist
  – Value exists but is unknown (n/a, not available)
  – Value not applicable

• The schema specifies for each attribute if it can be null (nullable attribute) or not with NOT NULL
NULL Values and Operators

For numerical operations:
- If x = NULL then \(4 \times (3-x)/7\) is also NULL

For boolean operations, in SQL there are three values:

<table>
<thead>
<tr>
<th>Value</th>
<th>SQL Value</th>
</tr>
</thead>
<tbody>
<tr>
<td>FALSE</td>
<td>0</td>
</tr>
<tr>
<td>TRUE</td>
<td>1</td>
</tr>
<tr>
<td>UNKNOWN</td>
<td></td>
</tr>
</tbody>
</table>

If x = NULL then x='Joe' is UNKNOWN

Note: comparison in SQL is a single '='

SQLite does not have a boolean datatype. It uses Integer instead!
Try:
- SELECT 2>1
- SELECT 2>NULL
- SELECT 1+NULL
Null Values in the WHERE Clause

```
SELECT * 
FROM   Person 
WHERE  (age < 25) 
   AND (height > 6 AND weight > 190) 
```

Will not return age=20, height=NULL, weight=200
Since NULL > 6 is UNKNOWN!
NULL Values in WHERE Clauses

Unexpected behavior:

```
SELECT * 
FROM Person 
WHERE age < 25 OR age >= 25
```

Should return all persons, but persons with NULL as age are not included!

You can use CASE with IS NULL, ISNULL(), IFNULL() or COALESCE() to handle NULL values.
CASTing Data Types

SQL is a typed language. I.e., values and columns have a data type.

```
SELECT 3/2
SELECT 3.0/2
SELECT 3/2.0
SELECT CAST(3 AS DOUBLE)/2
```

Typecasting rules are similar to other typed languages like C++. 
**RECAP: Inner Joins**

**Inner joins** select all rows from both tables as long as there is a match between the columns in both tables. Inner joins are the default in SQL.

**Example:** What stores sell what products?

```
SELECT Product.name, Purchase.store
FROM Product
JOIN Purchase ON Product.name = Purchase.prodName
```

Both equivalent:
```
SELECT Product.name, Purchase.store
FROM Product, Purchase
WHERE Product.name = Purchase.prodName
```

Both INNER JOINS!
Inner Joins + NULLS = Lost data?

However: Products that were never sold in any store (with no Purchase tuple) will be lost!
Outer Joins

An outer join returns also tuples from the joined relations that do not have a corresponding tuple in the other relations (filled with NULL values).

Left outer joins in SQL:

```sql
SELECT Product.name, Purchase.store
FROM Product
    LEFT OUTER JOIN Purchase
    ON Product.name = Purchase.prodName
```

Now we’ll get products even if they didn’t sell
INNER JOIN:

**Product**

<table>
<thead>
<tr>
<th>name</th>
<th>category</th>
</tr>
</thead>
<tbody>
<tr>
<td>Gizmo</td>
<td>gadget</td>
</tr>
<tr>
<td>Camera</td>
<td>Photo</td>
</tr>
<tr>
<td>OneClick</td>
<td>Photo</td>
</tr>
</tbody>
</table>

**Purchase**

<table>
<thead>
<tr>
<th>prodName</th>
<th>store</th>
</tr>
</thead>
<tbody>
<tr>
<td>Gizmo</td>
<td>Wiz</td>
</tr>
<tr>
<td>Camera</td>
<td>Ritz</td>
</tr>
<tr>
<td>Camera</td>
<td>Wiz</td>
</tr>
</tbody>
</table>

**SQL Query**

```
SELECT Product.name, Purchase.store
FROM Product
INNER JOIN Purchase
ON Product.name = Purchase.prodName
```
**LEFT OUTER JOIN:**

**Product**

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**Purchase**

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</table>

**SELECT** `Product.name, Purchase.store`

**FROM** `Product` **LEFT OUTER JOIN** `Purchase` **ON** `Product.name = Purchase.prodName`
Other Outer Joins

- **Left outer join:**
  - Include the left tuple even if there’s no match
- **Right outer join:**
  - Include the right tuple even if there’s no match
- **Full outer join:**
  - Include both left and right tuples even if there’s no match

SQLite currently only supports LEFT OUTER JOIN, but you can easily just change the order of the tables in the query.
Adding Data

INSERT INTO TABLE_NAME
    [(column1, column2, column3,...columnN)]
VALUES (value1, value2, value3,...valueN);

Note: column names are optional.

INSERT INTO Product
VALUES ('Gizmo', 19, 'Gadgets', 'GWorks')
Adding Data

The data can also come from an existing table.

```
INSERT INTO first_table_name [(column1, column2, ... columnN)]
  SELECT column1, column2, ...columnN
FROM second_table_name
WHERE condition;
```
Removing a Table

DROP TABLE database_name.table_name
Select Syntax Diagram (SQLite)

http://www.sqlite.org/lang.html
Activity


- Transaction control
- Views
- Indexes
- Date & Time